

CS W3134: Data Structures in Java

Lecture #7: Ordered lists, complexity, sorts

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Administrivia

- HW#1 due Thursday!
 - Any general questions?
- Akash has replaced William as TA
 - Note changed office hours

Agenda

- Finish ordered lists
- Big-Oh notation (complexity)
- Sorting algorithms
- Start stacks, if time allows

Ordered lists

- Find(), redux
 - What's the *stopping condition* for find()?

Costs

- How much do each of the previous operations cost in the *worst case*?
 - Most are linear, some are unit
- Binary search is special – it's better than linear time
 - Divide the range by half until too small to divide further == # of comparisons needed
 - Reverse: what's the range that can be covered with n steps? (Book page 63)
 - i.e., $r = 2^n$
 - What's this expressed as in terms of s ?
 - $s = \log_2 r$
 - Algorithm grows *logarithmically*

Formalizing costs

- We're going to approach this informally
- Time to insert one element is some constant K
 - e.g., $T(N) = K$
- Time to search for an element (linearly) is $T(N) = K * N$
- "Big-Oh Notation": upper-bound on worst-case time
 - We drop the constant K – for *sufficiently large* N , the constant is unimportant
 - To be precise, we find a function $F(x)$, where $T(x)$ is $O(F(x))$ if $|T(x)| \leq K|F(x)|$ for some $x > c$
 - The idea of doubling your computer's speed is embedded in K
 - $T(N) = O(N)$, for example

Examples of costs

- For lists using arrays?
 - Linear search: $O(N)$
 - Etc.
 - Draw a graph of the comparative costs, page 72
- What are bad about arrays?
 - Slow search in unordered, slow insert in ordered – can we speed both? Yes
 - Fixed size: can we change that? Yes

Sorts

- Bubble (p. 85)
 - Sort pairwise repeatedly
 - Biggest placed each time
- Selection (p. 89)
 - Search for smallest, swap with first
 - Search for smallest, swap with second
- Insertion (p. 95)
 - Take the next one, and put it into the existing sorted subset
- All $O(n^2)$
 - But they're not the exact same performance
- Let's write out a little bit of pseudocode for each

Sorts II

- Lexicographical comparisons?
- Stability of existing items?
- Sidebar: Comparable interface
 - All you have to do is implement boolean `compareTo(Object o)`
 - Generally a good thing to program to, I prefer to book's example
 - `Arrays.sort()`

Stacks and Queues

- Useful programmer's tools, will encounter it in many places
 - Very easy and fast to implement
 - Runs very fast as well
- "Restricted access": no index – only manipulate one item at a time
- More abstract – the underlying implementation is unimportant or not remotely similar to the structure, unlike lists

Stacks

- Basic operations: "LIFO" strategy
 - Push
 - Pop
 - Peek
- Analogy: mail basket
 - Not as rigorous as a real stack, of course
- Another analogy: life
 - Conversations
 - Workday
- Extraordinarily simple!

Array-based stacks

- Limited size; ways to get around this
- Decoupled from array index!
- Very simple to implement
 - Keep *top* variable, initialized to -1
- Boundary conditions?
- Complexity bounds?
 - Apart from simplicity, biggest reason to use

Next time...

- Reasons to use stacks
- Queues
- Arithmetic expression parsing
