## CS W3134: Data Structures in Java

Lecture \#15: Sorts
10/28/04
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## Administrivia

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- No class on Tuesday - vote!
- HW\#2 grades up Tuesday, can pick up on Thursday
- HW\#3 due today
- HW\#4 out tonight


## Agenda

- Three major types of fast sorts $\qquad$
- Mergesort
- Radix sort
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- Quicksort $\qquad$
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## Mergesort

Classic recursive algorithm

- Split arrays in half, sort each half, and then
$\qquad$ merge them together
- "Divide and conquer" $\qquad$
- Sort is the "recursive" call
- Psuedocode?
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## Mergesort (II)

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- Key aspect of code on page 287 $\qquad$
- The header of the method contains enough information to perform the recursive call $\qquad$
- In this case, partition information
- Efficiency?
- Partition: $\mathrm{O}(1)$
- Merge: $\mathrm{O}(\mathrm{n})$
- How many times each have to be done? $\mathrm{O}(\log \mathrm{n})$ $\qquad$
- Ergo, $\mathrm{O}(\mathrm{n} * \log \mathrm{n})$
- Disadvantage: lots of memory required


## Radix Sort

- Radix is the "base" of a system of numbers
- Very simple, fast algorithm (but a little tricky to implement)
- Sort by digit, one at a time
- Sort on the 1 s digit
- Sort on the 10 s digit; keep relative order of equal 10 s the same, i.e., go left-to-right on the 1 s digit
- Sort the 100 s digit
- Etc.
- Problem: where to store intermediate results?
- Can sort 100 numbers in 2 passes! $\sim \mathrm{O}(2 \mathrm{n})$
- But... that's essentially $O(n \log n)$ !
- There's no free lunch, but this works very well for specialized keys


## Quicksort: Partition

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- Relies on concept of partition
- A number s.t. two groups are formed: those smaller than the number, and those larger than the number
- "Pivot"
- Walk from both edges - If left is smaller than pivot, walk left
- If right is larger than pivot, walk right
- Otherwise, swap the two
- What if we cross?
- Last element is the pivot?
- Code? p. 338
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## Quicksort: Recursion

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- Given pivot, we:
- Partition the array in two;
- Quicksort the left "half";
- Quicksort the right "half".
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- And recurse!
- That's it (p. 338)
- Well, must be very, very careful
- Analysis?
- Usually $\mathrm{O}(\mathrm{n} \log \mathrm{n})$, and in-memory
- But there are some problems...


## Quicksort: Picking the pivot

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- Imagine a reverse-sorted array $\qquad$
- How long does Quicksort take? $\mathrm{O}\left(\mathrm{n}^{2}\right)$ !
- How can we fix this? $\qquad$
- Pick pivot more intelligently
- Two popular mechanisms: $\qquad$
- Random
- Median-of-three
- Also, inefficient for small arrays
- Use insertion sort as a degenerate case...


## Trees

- Linked Lists are generally connected to one other link
- What if we connect to multiple other links?
- A Tree is one generalization of a Linked List
- Key definition: no "cycles" amongst children - Graphs are more general
- Terminology
- Node, Edge, Path, Root, Parent, Child, Leaf, Subtree, Level
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| - Start trees | Next time |
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