

CS W3134: Data Structures in Java

Lecture #17: Quicksort, Trees

11/9/04

Janak J Parekh

Administrivia

- HW#4 due Thursday; last-minute questions?
 - Reminder: use the webboard to your advantage
- Important office hour changes
 - Akash will not hold office hours on Wednesday
 - Rachel will hold hers from 10:45am-12:45pm
 - Matthew will hold an extra hour from 6:30pm-7:30pm

Agenda

- Finish Quicksort
- Trees

Quicksort: Picking the pivot

- Imagine a reverse-sorted array
- How long does Quicksort take now? $O(n^2)$
- How can we fix this?
 - Pick pivot more intelligently
 - Two popular mechanisms:
 - Random
 - Median-of-three
- Also, inefficient for small arrays
 - Use insertion sort as a degenerate case...

Trees

- Linked Lists are generally connected to *one* other link
- What if we connect to multiple other links?
- A Tree is one generalization of a Linked List
- Key definition: no “cycles” amongst children
 - Graphs are more general
- Terminology
 - Node, Edge, Path, Root, Parent, Child, Leaf, Subtree, Level

Binary search trees

- What's a binary tree?
 - Two children, always
- Main concept:
 - Max(left subtree) must be $<$ current node, min(right subtree) must be $>$ current node
- Why?
 - Combines advantages of a linked list and an ordered array
 - Can insert fast and search fast
 - Unlimited growth
 - Relatively fast indexed access

Writing the Tree in Java

- “Node” class, with left and right children
- Data in node as well
- Very similar to Link
- Main “Tree” class that links to *root*, with find, insert, delete, etc. methods

Operations in a BST

- Search
 - Simple: walk left or right depending if $<$ or $>$ than current
 - If we hit the bottom, we can't find it
 - $O(\log N)$ time
- Insert
 - “Search”, and then put in the appropriate place
 - Need a “current” and a “parent” pointer, similar to linked-list

Traversing the tree

- Unlike search, want to walk in an abstract order, sort of like arrays
- Three means of traversal; all recursive
 - Inorder
 - Visit left subtree
 - Visit node
 - Visit right subtree
 - Preorder
 - Postorder
 - The latter two have use in expressions (pg. 386)

Other operations

- Min/max values
- Deleting a node
 - More complicated!
 - If no children, then nuke
 - One child
 - More than one child
 - Make one left, and go all the way right, or;
 - Make one right, and go all the way left
 - Take that node and put it at the deleted node's location
 - Move the right child of the moved node up one notch
 - Book uses latter convention

Tree complexity

- # of levels of a full tree is $\log N$
 - Search, insert, delete is $O(\log N)$
- What if it isn't full? Difficult analysis
 - Insert(1)
 - Insert(2)
 - ...
 - In fact, this is the one downside of simple BST trees: easy to make unbalanced
 - There are alternatives; you can read chapter 9 (optional)

Trees as arrays

- Array[0] is the root
- $2 * \text{index} + 1$ is the left child
- $2 * \text{index} + 2$ is the right child
- Parent of a node is, correspondingly, $(\text{index} - 1) / 2$
- Actually works surprisingly well, but...
 - No unlimited growth
 - Inefficient use of memory
 - Deletes are slow

Expression trees

- Operators are root and intermediate nodes, operands are leaf nodes
- To create
 - Start with postfix expression and a stack
 - Operand: form unit tree with value and push onto the stack
 - Operator: pop two things off of stack, combine “by” operator, push result on stack
- When done, one element on stack
- What does inorder, preorder, postorder mean?

Next time

- Finish trees
- Start hashing
