## **Nonogram Solver**

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### **1. Introduction**

## **1.1 Background:**



**Figure 1. Example nonogram puzzle**

In this project, we seek to parallelize a nonogram solver in Haskell. A nonogram is a visual puzzle where, given a grid of white cells, row constraints, and column constraints, players must fill out each cell to construct a picture. Row and column constraints are given for each row and column. "Blocks" are consecutively filled cells. Each constraint describes the number of blocks within a row or column and the lengths of each. From these constraints, players must iteratively infer which cells must be filled. Moreover, each block must be separated by at least one white or unfilled cell. Therefore, a row constraint like "3 2" describes a length 3 block, separated by at least one white cell, and followed by a length 2 block.

## **1.2 Challenges:**

The nonogram becomes challenging with difficult clues. For example, a row with less cells and with larger blocks is easier to solve. Conversely a row that is mostly composed of white cells introduces more ambiguity. Moreover, Puzzles designed for humans generally have one solution and reveal an image. However, it is also possible for Nonograms to have multiple solutions with no discernible pictures.

The nonogram also becomes more computationally complex with large grid sizes. As the number of rows and columns increases, the search space grows exponentially as there are more possibilities to explore. Nonograms have also been shown to be NP-hard and thus there are a

multitude of possible approaches to solving one that balance correctness with computational complexity.

## **2. Implementation**

## **2.1 Nonogram Representation:**

In our project, each nonogram is represented with key attributes of **height**, **width**, **rowArgs** and **colArgs**. **height** and **width** refer to the size of the grid of cells, while **rowArgs** and **colArgs** are lists of lists containing the constraints for each row and column. These are given as inputs to the algorithm to start solving the nonogram.

## **2.2 Data Collection and Parsing**

Nonograms were taken from [https://github.com/mikix/nonogram-db.](https://github.com/mikix/nonogram-db) We stored each nonogram from this database as a .txt file and used our **parseNonogram** function to extract rowArgs and colArgs (constraints) to be inputted into our solver.

## **2.3 Base Algorithm Overview:**

Our nonogram solver base algorithm can be described in three parts: 1) constraint satisfaction, 2) iterative inference, and 3) backtracking for unresolved cases.

# **1) Constraint Satisfaction**

The algorithm first iterates through each row and column constraint. For each, it comes up with possible placements of blocks. **cmputeBlocksSeq** :: Int -> [Int] -> [[Int]] operates on a line and takes the line length and constraints (block lengths) as input and recursively tries to place blocks into different start positions and continues with the remaining blocks. The output is a list of lists where each inner list contains possible start positions for the blocks within that line. For example, **lineLength** =  $7$  and **lineConstraint** =  $\begin{bmatrix} 2, 3 \end{bmatrix}$ would yield an output of [[0, 3], [1, 4]] from **computeBlocksSeq**.

```
computeBlocksSeq :: Int -> [Int] -> [[Int]]
computeBlocksSeq lineLength blockLengths = placeBlocks blockLengths 0
where
   -- Recursive helper function to place blocks
  placeBlocks :: [Int] -> Int -> [[Int]]
  placeBlocks [] = [[]] -- No blocks left to place
  placeBlocks (b:bs) start
     | start + remainingLength > lineLength = []
     | otherwise = do
        pos <- [start .. lineLength - remainingLength]
```

```
rest \leftarrow placeBlocks bs (pos + b + 1) -- Recur with updated start position
    return (pos : rest)
where
  -- Calculate remaining length
  remainingLength = sum (b : bs) + length bs
```
**generateBlocksSeq** :: [[Int]] -> [Int] -> Int -> [[Int]] takes the output of computeBlocks—the potential starting positions of the blocks—and generates a binary array (Ints of 1s and 0s) to represent possible line configurations.

```
generateBlocksSeq :: [[Int]] -> [Int] -> Int -> [[Int]]
generateBlocksSeq blockStarts blockSizes totalLength =
  map (generateBinaryArray blockSizes totalLength) blockStarts
where
  generateBinaryArray :: [Int] -> Int -> [Int] -> [Int]
  generateBinaryArray sizes len starts = foldl placeBlock (replicate len 0) (zip starts
sizes)
  placeBlock :: [Int] -> (Int, Int) -> [Int]
  placeBlock arr (start, size) =
      take start arr ++ replicate size 1 ++ drop (start + size) arr
```
### **2) Iterative inference**

With the possible line configurations, we then move on to the inference step. If in all possible configurations of a line, a cell is filled, then we know that that cell must be black. Similarly, if in all possible configurations, a cell is unfilled, then we know that that cell is certainly white. The function inferValues performs this step and is described as follows: **inferValuesSeq** :: PartialSolution -> (Array Int (Set.Set [Int]), Array Int (Set.Set [Int])) -> PartialSolution. A PartialSolution is a grid representing the current progress of the puzzle. This array contains values 0, 1, and -1. 0 represents white cells, 1 represents black cells, and -1 represents cells that are still unknown. **inferValuesSeq** is called repeatedly until the nonogram is solved.

The main logic in **inferValuesSeq** is in the following code from the function, and also in a helper function **inferRowOrCol**:

```
let rowOnes = foldl1 (zipWith (.&.)) placements
   rowZeros = foldl1 (zipWith (.|.)) placements
   inferredRow = inferRowOrCol rowOnes rowZeros
in [((r, c), inferredRow !! c) | c \leftarrow [0..numCols - 1], partialSolution !(r, c) == -1]-- Infer the result for a row or column based on bitwise results
```

```
inferRowOrCol :: [Int] -> [Int] -> [Int]
```

```
inferRowOrCol ones zeros =
  zipWith resolveCell ones zeros
where
  resolveCell 0 1 = -1 -- The cell must be unknown
   resolved 1 1 1 = 1 - - The cell must be filled (from ones)
   resolved100 = 0 -- The cell must be empty (from zeros)
  resolved 1 = -1 - -  0.5 Fallback for unexpected values
```
We perform bitwise operations to extract indices where it is 1s or 0s across all possible placements for that line.

## **3) Backtracking**

If the nonogram puzzle is ambiguous—meaning that there are multiple solutions and options which cannot be deduced by the iterative solver alone—then we use backtracking to solve the rest of the puzzle. Backtracking works by trying different placements within a row and checking if this step still results in a valid grid. If the grid is still valid, backtrack is called recursively until we reach the basecase of all the rows being completed. In our implementation, backtrack is defined as follows:

```
backtrack :: PartialSolution
       -> Array Int (Set [Int]) -- Row placements
       -> [Constraint] -- Row constraints
       -> [Constraint] -- Column constraints
       -> Set Int -- Completed rows
       -> Set Int -- Completed columns
       -> [PartialSolution] -- Accumulated solutions
       -> [PartialSolution] -- Final list of valid solutions
```
### **3. Sequential Algorithm Benchmark**





**Table 1. Sequential algorithm benchmark with puzzles of different sizes**

**Figure 3. Activity graph for sequential algorithm with 7 puzzles**

We first tested our sequential algorithm on puzzles of varying difficulty. As expected, smaller puzzles were solved very quickly (0.01s), and larger puzzles took more time as there is a larger search space for the algorithm to go through. As one puzzle may be too little work, we decided on the benchmark task being to solve 7 medium puzzles. Larger puzzles with sizes 60x70 were also considered and tested, however the terminal would crash with these sizes, so we settled on testing the algorithm on multiple puzzles instead.

# **4. Parallelization Strategy**

Our motivation for parallelization came from the fact that row and column processing can be done independently. For instance, computing the starting placements based on the constraints for one row is not affected by the result of another row. The Control.Parallel.Strategies module was used for parallelization and our attempts target separate functions that contribute to the main algorithm.

# **4.1 Parallelization Attempt 1: inferValuesPar**

We first attempted to parallelize **inferValues** as each inference for each line is done independently. Specifically, **parMap** and **rseq** were used to allow multiple rows to be processed simultaneously with **updateRows**. inferValues is mainly used in iterativeSolve which calls the function repeatedly. The iterativeSolve function is inherently sequential, as the partial solution grid must be updated before another call can be made in the next step. Because of this, we do not expect significant gains from this parallelization strategy. Below is the line of code where parallelization occurs.

**updatedRows** = partialSolution // concat (parMap rseq updateRow [r1..r2])

#### **4.2 Parallelization Attempt 2: computeBlocksPar**

The second parallelization attempt involves parallelizing computeBlocks—extracting valid starting positions for each block from the constraints. To be able to control the granularity of the task, we chunked multiple cells to be processed together using a helper function chunkList. We noticed that computeBlocksPar benefits from doing this. processChunk applies processPosition which computes the valid placements for a single starting position. parMap and rdeepseq were used to ensure parallelization across threads.

```
computeBlocksPar lineLength blockLengths =
   concat $ parMap rdeepseq processChunk (chunkList chunkSize [0 .. lineLength -
totalRemainingLength])
```
### **4.3 Parallelization Attempt 3: generateBlocksPar**

generateBlocksPar uses the same method from computeBlocksPar for parallelization. It groups the cells in chunks using chunkList. processChunk then applies generateBinaryArray which performs the main task of placing the blocks in the binary array. parMap and rdeepseq were used for parallelization.

**generateBlocksPar** blockStarts blockSizes totalLength = concat \$ parMap rdeepseq processChunk (chunkList chunkSize blockStarts)

#### **4.4 Parallelization Summary**

With these attempts, we accumulate two different versions (sequential and parallel) of each inferValue, computeBlock, and generateBlocks. For each, you can choose whether to use the sequential version or the parallelized version. We test different combinations of the three functions. For example **solveParallelComputeGenerate** uses both the parallelized version of computeBlocks and generateBlocks, while **solveFullyParallel** uses the parallelized version of computeBlocks, generateBlocks, and inferValues. In total, there are 6 versions of the algorithm:

```
solveSequential :: FilePath -> IO ()
solveSequential = solveNonogramFromFile computeBlocksSeq generateBlocksSeq iterativeSolveSeq
solveParallelComputeBlocks :: FilePath -> IO ()
solveParallelComputeBlocks = solveNonogramFromFile computeBlocksPar generateBlocksSeq iterativeSolveSeq
solveParallelGenerateBlocks :: FilePath -> IO ()
solveParallelGenerateBlocks = solveNonogramFromFile computeBlocksSeq generateBlocksPar iterativeSolveSeq
solveParallelComputeGenerate :: FilePath -> IO ()
solveParallelComputeGenerate = solveNonogramFromFile computeBlocksPar generateBlocksPar
iterativeSolveSeq
```

```
solveParallelIterativeSolve :: FilePath -> IO ()
solveParallelIterativeSolve = solveNonogramFromFile computeBlocksSeq generateBlocksSeq iterativeSolvePar
solveFullyParallel :: FilePath -> IO ()
solveFullyParallel = solveNonogramFromFile computeBlocksPar generateBlocksPar iterativeSolvePar
```
## **4. Results**



### **Table 2. Speedup graph with different algorithms**

The table above shows the speedup, which was calculated as (Sequential benchmark) / (shortest time elapsed for algorithm). Parallel Generate performed the best with a x1.88 speed up from the original benchmark, followed by Parallel Iterative Solve with a x1.50 improvement. Fully Parallel and Parallel Compute Generate did not perform as well. This shows that combining multiple parallelized functions is detrimental to speed up time, and causes more overhead computations than it is worth.



**Figure 4. Total time elapsed vs. number of threads**

The graph above shows the total time elapsed with increasing number of threads. For all parallelized versions of the algorithm, the time elapsed decreased, showing utilization of the threads. However, all of the algorithms also level out, showing diminishing returns from an excessive number of threads.

## **5. Improvements**



### **5.1 Limitation from Iterative Solve**

## **Figure 5. Activity graph for Fully Parallel -N8**

As shown in the activity graph (Figure 5), we struggled to parallelize the main process in the algorithm — iterativeSolve, which repeatedly infers values until the nonogram is solved. This is because the nature of iterativeSolve is sequential, where the next step depends on the result of the previous. However, the activity graph does show somewhat successful parallelization when doing computeBlocksPar and generateBlocksPar.

As the puzzles were not ambiguous, backtracking did not end up being used. For the next steps, however, backtracking part of the algorithm would likely gain more from parallelization as multiple placements can be explored concurrently. We can focus on solving smaller nonograms, but mainly through backtracking.



# **5.2 Garbage Collection**

**Figure 6. GC Time vs. MUT Time for Parallel Generate -N8**

We also noticed that garbage collection takes up as much time as mutator operations (Figure 6). This means that a significant portion of the time is spent on managing memory rather than doing useful work. A future solution could be to try to reduce GC time, perhaps using ParBuffer, which could help limit memory usage.

# **6. References**

- <https://github.com/mikix/nonogram-db>
- <https://github.com/Arpanio/nonogram>
- [https://towardsdatascience.com/solving-nonograms-with-120-lines-of-code-a7c6e0f627e](https://towardsdatascience.com/solving-nonograms-with-120-lines-of-code-a7c6e0f627e4) [4](https://towardsdatascience.com/solving-nonograms-with-120-lines-of-code-a7c6e0f627e4)

### **7. Code**

### **7.1 Main.hs**

```
import NonogramSolverPar (solveSequential, solveParallelComputeBlocks,
solveParallelGenerateBlocks, solveParallelComputeGenerate,
solveParallelIterativeSolve, solveFullyParallel)
import System.Environment (getArgs)
main :: IO ()
main = do
  putStrLn "Solving Nonogram Puzzle..."
  args <- getArgs
  case args of
      [mode, filePath] ->
          case mode of
               "--sequential" -> solveSequential filePath
               "--parallel-compute" -> solveParallelComputeBlocks filePath
               "--parallel-generate" -> solveParallelGenerateBlocks filePath
               "--parallel-compute-generate" -> solveParallelComputeGenerate filePath
               "--parallel-iterative-solve" -> solveParallelIterativeSolve filePath
               "--fully-parallel" -> solveFullyParallel filePath
               _ -> putStrLn "Invalid mode. Usage: ./nonogramSolver [--sequential |
 -parallel-compute | --parallel-generate | --parallel-compute-generate |
 -parallel-iterative-solve | --fully-parallel] <file-path>"
       -> putStrLn "Usage: ./nonogramSolver [--sequential | --parallel-compute |
 -parallel-generate | --parallel-compute-generate | --parallel-iterative-solve |
  putStrLn "Puzzle Solved!"
```
#### **7.2 Parser.hs**

```
module Parser (parseNonogram) where
import NonogramTypes
import Data.List (isPrefixOf, uncons)
```

```
import Data.List.Split (splitOn)
import Data.Maybe (mapMaybe)
 -- Parse
parseNonogram :: FilePath -> IO Nonogram
parseNonogram path = do
content <- lines <$> readFile path
let titleLine = extractValue "title" content
    heightLine = read (extractValue "height" content) :: Int
    widthLine = read (extractValue "width" content) :: Int
    rowsSection = extractSection "rows" "columns" content
    colsSection = extractSection "columns" "goal" content
    goalSection = extractValue "goal" content
    rowsHints = parseHints rowsSection
    colsHints = parseHints colsSection
 return Nonogram { title = titleLine, height = heightLine, width = widthLine, rows =
rowsHints, columns = colsHints, goal = goalSection }
extractValue :: String -> [String] -> String
extractValue key allLines =
let matchingLines = [line | line <- allLines, key `isPrefixOf` line]
in case uncons matchingLines of
      Just (matchingLine, \_) -> unquote $ drop (length key + 1) matchingLine
     Nothing -> error $ "Key not found: " ++ key -- empty case
where
  unquote s = filter ('notElem' "\"") s -- no quotes
 - Extract sections
extractSection :: String -> String -> [String] -> [String]
extractSection startKey endKey allLines =
takeWhile (not . isPrefixOf endKey) . drop 1 . dropWhile (not . isPrefixOf startKey)
$ allLines
-- Parse hints into list of lists of integers
parseHints :: [String] -> [[Int]]
parseHints = mapMaybe safeParseLine
```

```
safeParseLine :: String -> Maybe [Int]
safeParseLine line =
if null line then Nothing
else Just (map read $ splitOn "," line)
```
## **7.3 NonogramTypes.hs**



### **7.4 NonogramSolverPar.hs**

```
module NonogramSolverPar (solveSequential, solveParallelComputeBlocks,
solveParallelGenerateBlocks, solveParallelComputeGenerate,
solveParallelIterativeSolve, solveFullyParallel) where
import Data.Array
import Data.List (group, transpose, foldl')
import Data.Set (Set)
import qualified Data.Set as Set
import Data.Bits ((.&.), (.|.))
import Parser (parseNonogram)
import NonogramTypes
import Control.Parallel.Strategies (parMap, rdeepseq, rseq)
computeBlocksSeq :: Int -> [Int] -> [[Int]]
computeBlocksSeq lineLength blockLengths = placeBlocks blockLengths 0
```

```
where
  placeBlocks :: [Int] -> Int -> [[Int]]
  placeBlocks [] = [[]] -- No blocks left to place
  placeBlocks (b:bs) start
     | start + remainingLength > lineLength = []
     | otherwise = do
        pos <- [start .. lineLength - remainingLength]
        rest \leq placeBlocks bs (pos + b + 1)
        return (pos : rest)
    where
       remainingLength = sum (b : bs) + length bs
computeBlocksPar :: Int -> [Int] -> [[Int]]
computeBlocksPar lineLength blockLengths =
  concat $ parMap rdeepseq processChunk (chunkList chunkSize [0 .. lineLength -
totalRemainingLength])
where
   totalRemainingLength = sum blockLengths + length blockLengths - 1
  chunkSize = 10
  -- Divide list into chunks
  chunkList :: Int \rightarrow [a] \rightarrow [[a]]
   chunkList [] = []
   chunkList n xs = take n xs : chunkList n (drop n xs)
  processChunk :: [Int] -> [[Int]]
  processChunk positions = concatMap processPosition positions
  processPosition :: Int -> [[Int]]
  processPosition start = placeBlocks blockLengths start
  placeBlocks :: [Int] -> Int -> [[Int]]
  placeBlocks [] = [[[]] -- No blocks left to place
  placeBlocks (b:bs) start
     | start + remainingLength > lineLength = []
    | otherwise = do
```

```
pos <- [start .. lineLength - remainingLength]
         rest \leq placeBlocks bs (pos + b + 1)
         return (pos : rest)
    where
       remainingLength = sum (b : bs) + length bs
generateBlocksSeq :: [[Int]] -> [Int] -> Int -> [[Int]]
generateBlocksSeq blockStarts blockSizes totalLength =
  map (generateBinaryArray blockSizes totalLength) blockStarts
where
  generateBinaryArray :: [Int] -> Int -> [Int] -> [Int]
  generateBinaryArray sizes len starts = foldl placeBlock (replicate len 0) (zip
starts sizes)
  placeBlock :: [Int] -> (Int, Int) -> [Int]
  placeBlock arr (start, size) =
       take start arr ++ replicate size 1 ++ drop (start + size) arr
-- Parallelized generateBlocks with chunking
generateBlocksPar :: [[Int]] -> [Int] -> Int -> [[Int]]
generateBlocksPar blockStarts blockSizes totalLength =
  concat $ parMap rdeepseq processChunk (chunkList chunkSize blockStarts)
  chunkSize = 10
  chunkList :: Int \rightarrow [a] \rightarrow [[a]]
  chunkList _ [] = []
  chunkList n xs = take n xs : chunkList n (drop n xs)
  processChunk :: [[Int]] -> [[Int]]
  processChunk chunk = map (generateBinaryArray blockSizes totalLength) chunk
  generateBinaryArray :: [Int] -> Int -> [Int] -> [Int]
  generateBinaryArray sizes len starts =
       foldl' placeBlock (replicate len 0) (zip starts sizes)
```

```
placeBlock :: [Int] -> (Int, Int) -> [Int]
  placeBlock arr (start, size) =
       take start arr ++ replicate size 1 ++ drop (start + size) arr
validGroups :: [Int] -> String -> Bool
validGroups expected line = groupLengths == expected
where
  groupLengths = map length . filter (all (== '1')) $ group line
type Constraint = [Int]
type PartialSolution = Array (Int, Int) Int
-- Helper function: Extract groups of 1s from a row/column as strings
extractGroupsAsString :: [Int] -> String
extractGroupsAsString = map toChar
where
  toChar 1 = '1'toChar 0 = '0'toChar = '-'validateLine :: [Int] -> Constraint -> Bool
validateLine line constraint
 | all (== -1) line = True -- Skip untouched lines
 | sum (filter (== 1) line) > sum constraint = False -- Exceeding constraints
 | any (\{x \rightarrow x == -1\}) line && sum (filter (== 1) line) <= sum constraint = True --
Partially solved line
 | otherwise = validGroups constraint (extractGroupsAsString line)
valid :: [Constraint] -> [Constraint] -> PartialSolution -> Set.Set Int -> Set.Set Int
\Rightarrow Bool
valid rowArgs colArgs partialSolution completedRows completedCols =
let
```

```
((r1, c1), (r2, c2)) = bounds partialSolution
  localRows = [[partialSolution ! (r, c) | c <- [c1..c2]] | r <- [r1..r2]]
  cols = transpose localRows
  validRows = all (\n\langle r, row \rangle -\n\rangle)let
      constraint = rowArgs !! r
      if r `Set.member` completedRows
      then validGroups constraint (extractGroupsAsString row)
      else validateLine row constraint
    ) $ zip [0..] localRows
   -- Validate columns
  validCols = all (\setminus (c, col) ->
    let
      constraint = colArgs !! c
      if c `Set.member` completedCols
      then validGroups constraint (extractGroupsAsString col)
      else validateLine col constraint
    ) $ zip [0..] cols
 in
  validRows && validCols
inferValuesSeq :: PartialSolution -> (Array Int (Set.Set [Int]), Array Int (Set.Set
[Int])) -> PartialSolution
inferValuesSeq partialSolution (rowPlacements, colPlacements) =
   let ((r1, c1), (r2, c2)) = bounds partialSolution
      numCols = c2 - c1 + 1numRows = r2 - r1 + 1updatedRows = partialSolution // concat (parMap rseq updateRow [r1..r2])
        where
          updateRow r =let placements = Set.toList $ rowPlacements ! r
              in if null placements
```

```
else
                     let rowOnes = foldl1 (zipWith (. & .)) placements
                          rowZeros = foldl1 (zipWith (.|.)) placements
                          inferredRow = inferRowOrCol rowOnes rowZeros
                      in [((r, c), infercdRow!! c) | c < -[0..numCols - 1],partialSolution ! (r, c) = -1]
      updatedCols = updatedRows // concatMap updateCol [c1..c2]
        where
          updateCol c =let placements = Set.toList $ colPlacements ! c
              in if null placements
                 else
                     let colOnes = foldl1 (zipWith (.&.)) placements
                         colZeros = fold11 (zipWith (.|.)) placementsinferredCol = inferRowOrCol colOnes colZeros
                     in [((r, c), inferredCol !! r) |r < - [0.. numRows -1],
updatedRows ! (r, c) = -1]
  in updatedCols
inferValuesPar :: PartialSolution -> (Array Int (Set.Set [Int]), Array Int (Set.Set
[Int])) -> PartialSolution
inferValuesPar partialSolution (rowPlacements, colPlacements) =
  let ((r1, c1), (r2, c2)) = bounds partialSolution
      numCols = c2 - c1 + 1updatedRows = partialSolution // concat (parMap rseq updateRow [r1..r2])
        where
          updateRow r =let placements = Set.toList $ rowPlacements ! r
              in if null placements
                 else
                     let rowOnes = foldl1 (zipWith (. & .)) placements
                          rowZeros = foldl1 (zipWith (.|.)) placements
                         inferredRow = inferRowOrCol rowOnes rowZeros
```

```
in [(r, c), inferredRow !! c) | c < - [0..numCols - 1],
partialSolution ! (r, c) = -1]
      numRows = r2 - r1 + 1updatedCols = updatedRows // concatMap updateCol [c1..c2]
        where
          updateCol c =let placements = Set.toList \frac{1}{5} colPlacements ! c
              in if null placements
                 then [] -- If there are no placements, skip updates for this column
                 else
                      let colOnes = foldl1 (zipWith (.&.)) placements
                          colZeros = foldl1 (zipWith (.|.)) placements
                          inferredCol = inferRowOrCol colOnes colZeros
                      in [((r, c), inferredCol !! r) |r < - [0..numRows - 1],
updatedRows ! (r, c) == -1]
  in updatedCols
 -- Infer result for row or column based on bitwise results
inferRowOrCol :: [Int] -> [Int] -> [Int]
inferRowOrCol ones zeros =
  zipWith resolveCell ones zeros
where
  resolveCell 0 1 = -1 -- unknown
  resolveCell 1 1 = 1 - - filled (from ones)resolveCell 0 0 = 0 -- empty (from zeros)
  resolveCell = = -1 - - unexpected
-- Determine completed localRows and columns in solution array
-- Output: A tuple of two sets. localRows/colCompleted: indices of localRows/cols that
are completed,
updateCompletions :: Array (Int, Int) Int -> (Set Int, Set Int)
updateCompletions solutionArray =
  let ((r1, c1), (r2, c2)) = bounds solutionArray
      -- Extract localRows and check if all values in each row are not -1
```

```
rowsCompleted = Set.fromList [r | r < - [r1..r2], all ( /= -1) [solutionArray !
(r, c) | c \leftarrow [c1..c2]]-- Extract columns and check if all values in each column are not -1
       colsCompleted = Set.fromList [c | c \leftarrow [c1..c2], all | (-1) [solutionArray !
(r, c) | r \leftarrow [r1..r2]]in (rowsCompleted, colsCompleted)
type Placement = [Int]
type PlacementsDict = (Array Int (Set.Set [Int]), Array Int (Set.Set [Int])) --
(rowPlacements, colPlacements)
(updated rowPlacements & colPlacements)
 -- Calls: updateRow, isValidRow, updateCol, isValidCol
updatePlacements :: Array (Int, Int) Int -> PlacementsDict -> Set.Set Int -> Set.Set
Int -> (Bool, PlacementsDict)
updatePlacements solutionArray (rowPlacements, colPlacements) completedRows
completedCols =
  let
       -- Update localRows
       (updatedRows, updatedRowPlacements) = processRows solutionArray rowPlacements
completedRows
       -- Update columns
       (updatedCols, updatedColPlacements) = processColumns solutionArray
colPlacements completedCols
       -- Combine results
      updated = updatedRows || updatedCols
       (updated, (updatedRowPlacements, updatedColPlacements))
processRows :: Array (Int, Int) Int -> Array Int (Set.Set Placement) -> Set.Set Int ->
(Bool, Array Int (Set.Set Placement))
processRows solutionArray rowPlacements completedRows =
   let
       indexesForDeletion = [i | i < - indices rowPlacements, i `Set.member`
completedRows]
       validatePlacement rowIdx placement = all isValidCell (zip [0..] placement)
           where
               isValidCell (cellIdx, value) =
```

```
(value /= 0 || solutionArray ! (rowIdx, cellIdx) /= 1) &&
                   (value /= 1 || solutionArray ! (rowIdx, cellIdx) /= 0)
       processRow (rowUpdated, acc) rowIdx
           | rowIdx `Set.member` completedRows = (rowUpdated, acc // [(rowIdx,
Set.empty)])
           | otherwise =
                   validPlacements = Set.filter (validatePlacement rowIdx)
(rowPlacements ! rowIdx)
               in
                   if validPlacements /= rowPlacements ! rowIdx
                       then (True, acc // [ (rowIdx, validPlacements)]) -- make updates
to acc array
                       else (rowUpdated, acc)
       (updated, newPlacements) = foldl' processRow (False, rowPlacements) (indices
rowPlacements)
  in
       (updated, newPlacements)
-- Process columns: remove completed columns and validate remaining placements
processColumns :: Array (Int, Int) Int -> Array Int (Set.Set Placement) -> Set.Set Int
-> (Bool, Array Int (Set.Set Placement))
processColumns solutionArray colPlacements completedCols =
   let
       indexesForDeletion = [i \mid i \le -i indices colPlacements, i `Set.member`
completedCols] -- create a list of indices to remove later
       validatePlacement colIdx placement = all isValidCell (zip [0..] placement) --
           where -- iterates over each cell in placement to get row indx and value
               isValidCell (cellIdx, value) =
                   (value /= 0 || solutionArray ! (cellIdx, colIdx) /= 1) &&
                   (value /= 1 || solutionArray ! (cellIdx, colIdx) /= 0)
      processCol (colUpdated, acc) colIdx
           | colIdx `Set.member` completedCols = (colUpdated, acc // [(colIdx,
Set.empty)]) -- mark column as empty if completed
           | otherwise =
              let
                   validPlacements = Set.filter (validatePlacement colIdx)
(colPlacements ! colIdx)
                   if validPlacements /= colPlacements ! colIdx
```

```
then (True, acc // [(colIdx, validPlacements)])
                      else (colUpdated, acc)
      (updated, newPlacements) = foldl' processCol (False, colPlacements) (indices
colPlacements)
      (updated, newPlacements)
backtrack :: PartialSolution
        -> [Constraint] -- Row constraints
        -> Set Int -- Completed localRows
        -> Set Int -- Completed columns
backtrack solutionArray rowPlacements rowArgs colArgs completedRows completedCols
solutions
  | Set.size completedRows == length rowArgs =
      if not (any (== solutionArray) solutions)
      then solutionArray : solutions
      else solutions
  | otherwise =
      -- Iterate over localRows in rowPlacements
      foldl tryPlacement solutions (assocs rowPlacements)
where
  tryPlacement :: [PartialSolution] -> (Int, Set.Set [Int]) -> [PartialSolution]
  tryPlacement solList (row, options) =
      if Set.member row completedRows
      then solList -- Skip completed localRows
      else foldl (tryOption row) solList (Set.toList options)
  tryOption :: Int -> [PartialSolution] -> [Int] -> [PartialSolution]
  tryOption row solList option =
      let
```

```
originalRow = [solutionArray ! (row, col) | col <- colIndices]
          updatedSolution = solutionArray // [((row, col), option !! (col - cl)) ]col <- colIndices]
          (completedRowsNext, completedColsNext) = updateCompletions updatedSolution
          newRowPlacements = rowPlacements // [(row, Set.empty)]
      if not (valid rowArgs colArgs updatedSolution completedRowsNext
completedColsNext)
      then solList
      else
          -- Append solutions returned from recursive call to current list
          backtrack updatedSolution newRowPlacements rowArgs colArgs
completedRowsNext completedColsNext solList
      where
          ((_r \ c1), (_r \ c2)) = bounds solutionArray
          collindices = [c1..c2]iterativeSolveSeq :: PartialSolution
             -> [Constraint] -- Row constraints
             -> [Constraint] -- Column constraints
             -> Set Int -- Completed columns
iterativeSolveSeq solutionArray (rowPlacements, colPlacements) rowArgs colArgs
completedRows completedCols =
      inferredSolution = inferValuesSeq solutionArray (rowPlacements, colPlacements)
      (newCompletedRows, newCompletedCols) = updateCompletions inferredSolution
      -- Update placements
       (updatedFlag, newPlacements) = updatePlacements inferredSolution
(rowPlacements, colPlacements) newCompletedRows newCompletedCols
       (newRowPlacements, newColPlacements) = newPlacements
  in
      if not updatedFlag
      then (inferredSolution, (newRowPlacements, newColPlacements), newCompletedRows,
newCompletedCols)
```

```
else iterativeSolveSeq inferredSolution (newRowPlacements, newColPlacements)
rowArgs colArgs newCompletedRows newCompletedCols
iterativeSolvePar :: PartialSolution
             -> Set Int -- Completed localRows
iterativeSolvePar solutionArray (rowPlacements, colPlacements) rowArgs colArgs
completedRows completedCols =
  let
      inferredSolution = inferValuesPar solutionArray (rowPlacements, colPlacements)
      (newCompletedRows, newCompletedCols) = updateCompletions inferredSolution
       (updatedFlag, newPlacements) = updatePlacements inferredSolution
(rowPlacements, colPlacements) newCompletedRows newCompletedCols
       (newRowPlacements, newColPlacements) = newPlacements
      if not updatedFlag
       then (inferredSolution, (newRowPlacements, newColPlacements), newCompletedRows,
newCompletedCols)
      else iterativeSolvePar inferredSolution (newRowPlacements, newColPlacements)
rowArgs colArgs newCompletedRows newCompletedCols
printSolution :: PartialSolution -> IO ()
printSolution solution = do
  let ((r1, c1), (r2, c2)) = bounds solution
  mapM (putStrLn . concatMap show) [[solution ! (r, c) | c <- [c1..c2]] | r <-
[r1..r2]]
  putStrLn "" -- Add a blank line between solutions
```

```
solveNonogramFromFile :: (Int -> [Int] -> [[Int]]) --- computeBlocks
function
                    -> ([[Int]] -> [Int] -> Int -> [[Int]]) -- generateBlocks
[Constraint] -> Set Int -> Set Int
iterativeSolve function
                    -> FilePath -- File path to Nonogram
                    \Rightarrow IO ()
solveNonogramFromFile computeBlocksFunc generateBlocksFunc iterativeSolveFunc filePath
= do
  nonogram <- parseNonogram filePath
  -- Extract row and column arguments
  let rowArgs = rows nonogram
      colArgs = columns nonogram
  putStrLn $ "\nTitle: " ++ title nonogram
  mapM_ print rowArgs
  -- Compute row and column vector lengths
  let rowVectorLen = length colArgs
  let colVectorLen = length rowArgs
  let rowPlacements = listArray (0, length rowArgs - 1) $map (\arg -> Set.fromList (generateBlocksFunc (computeBlocksFunc
rowVectorLen arg) arg rowVectorLen)) rowArgs
  let colPlacements = listArray (0, length colArgs - 1) $
          map (\arg -> Set.fromList (generateBlocksFunc (computeBlocksFunc
colVectorLen arg) arg colVectorLen)) colArgs
```

```
let partialSolution = array ((0, 0), (length rowArgs - 1, length colArgs - 1))[(r, c), -1) | r <- [0..length rowArgs - 1], c <- [0..lengthcolArgs - 1]]
  let completedRows = Set.empty
  let completedColumns = Set.empty
  let placementsDict = (rowPlacements, colPlacements)
   let (finalSolution, \overline{\phantom{a}}, \overline{\phantom{a}}, \overline{\phantom{a}}) =
           iterativeSolveFunc partialSolution placementsDict rowArgs colArgs
completedRows completedColumns
  printSolution finalSolution
solveNonogramBacktrack :: (Int -> [Int] -> [[Int]]) -- computeBlocks
function
                      -> ([[Int]] -> [Int] -> Int -> [[Int]]) -- generateBlocks
function
                      -> (PartialSolution -> Array Int (Set [Int]) -> [Constraint] ->
[Constraint] -> Set Int -> Set Int
                          -> [PartialSolution] -> [PartialSolution]) -- backtrack
function
Nonogram
                      \Rightarrow IO ()
solveNonogramBacktrack computeBlocksFunc generateBlocksFunc backtrackFunc filePath =
do
  nonogram <- parseNonogram filePath
   -- Extract row and column arguments
  let rowArgs = rows nonogram
       colArgs = columns nonogram
  putStrLn $ "\nTitle: " ++ title nonogram
```

```
let rowVectorLen = length colArgs
  let colVectorLen = length rowArgs
  let rowPlacements = listArray (0, length rowArgs - 1) $
          map (\arg -> Set.fromList (generateBlocksFunc (computeBlocksFunc
rowVectorLen arg) arg rowVectorLen)) rowArgs
  let colPlacements = listArray (0, length colArgs - 1) $
          map (\arg -> Set.fromList (generateBlocksFunc (computeBlocksFunc
colVectorLen arg) arg colVectorLen)) colArgs
  let partialSolution = array ((0, 0), (length rowArgs - 1, length colArgs - 1))
                      [(r, c), -1) | r < - [0..length rowArgs - 1], c < - [0..length
colArgs - 1]]
  let completedRows = Set.empty
  let completedColumns = Set.empty
  putStrLn "Solving via backtracking..."
  let solutions = backtrack partialSolution rowPlacements rowArgs colArgs
completedRows completedColumns []
  putStrLn "Found solutions:"
  mapM_ printSolution solutions
```
solveSequential :: FilePath -> IO ()

```
solveSequential = solveNonogramFromFile computeBlocksSeq generateBlocksSeq
iterativeSolveSeq
solveParallelComputeBlocks :: FilePath -> IO ()
solveParallelComputeBlocks = solveNonogramFromFile computeBlocksPar generateBlocksSeq
iterativeSolveSeq
solveParallelGenerateBlocks :: FilePath -> IO ()
solveParallelGenerateBlocks = solveNonogramFromFile computeBlocksSeq generateBlocksPar
iterativeSolveSeq
solveParallelComputeGenerate :: FilePath -> IO ()
solveParallelComputeGenerate = solveNonogramFromFile computeBlocksPar
generateBlocksPar iterativeSolveSeq
solveParallelIterativeSolve :: FilePath -> IO ()
solveParallelIterativeSolve = solveNonogramFromFile computeBlocksSeq generateBlocksSeq
iterativeSolvePar
solveFullyParallel :: FilePath -> IO ()
solveFullyParallel = solveNonogramFromFile computeBlocksPar generateBlocksPar
iterativeSolvePar
```
### **7.5 TestNonogramSolver.hs**

```
module Main (main) where
import Test.HUnit
import NonogramSolverPar (solveSequential)
import NonogramTypes (Nonogram(..))
import Parser (parseNonogram)
import Data.Array (Array, bounds, (!))
import System.IO.Silently (capture)
testPuzzle :: FilePath -> Test
testPuzzle puzzlePath = TestCase $ do
```

```
nonogram <- parseNonogram puzzlePath
  let expectedSolution = goal nonogram
   if null expectedSolution
       then assertFailure $ "No solution provided in puzzle file: " ++ puzzlePath
      else do
           (actualOutput, _) <- capture (solveSequential puzzlePath)
          let actualSolution = cleanOutput actualOutput
          assertEqual ("Mismatch for puzzle: " ++ puzzlePath) expectedSolution
actualSolution
-- Helper function to reformat captured output
cleanOutput :: String -> String
cleanOutput = concat . map (filter (`elem` "01")) . lines
-- test cases
tests :: Test
tests = TestList
   [ testPuzzle "test/test_cases/bloop.txt",
       testPuzzle "test/test cases/spade.txt" ]
-- Run the tests
main :: IO ()
main = do
  testResults <- runTestTT tests
  print testResults
```