Haskell Parallel Chess Engine

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Introduction

Minimax based chess engine

Bitboards

Parallelization

Live demo

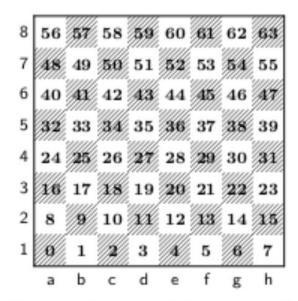


Figure 1: Chess board squares with the corresponding number

Bitboards

- Each square corresponds to a bit in a 64 bit word.
- We can easily map bits to squares
- Each piece can be shown as a 64 bit word

	а	b	с	d	e	f	g	h
1	0	1	2	3	4	5	6	7
2	8	9	10	11	12	13	14	15
3	16	17	18	19	20	21	22	23
4	24	25	26	27	28	29	30	31
5	32	33	34	35	36	37	38	39
6	40	41	42	43	44	45	46	47
7	48	49	50	51	52	53	54	55
8	56	57	58	59	60	61	62	63

Bitboards

startpos :: Board startpos = Board { pawnsWhite = 0x00000000000FF00 -- [a,h]2 entire row , pawnsBlack = 0x00FF00000000000 -- [a,h]7 entire row knightsWhite = 0x0000000000000042 -- b1 (bit 1) and g1 (bit 6) knightsBlack = 0x42000000000000 -- b8 (bit 57) and q8 (bit 62) bishopsWhite = 0x000000000000024 -- c1 (bit 2) and f1 (bit 5)bishopsBlack = 0x24000000000000 -- c8 (bit 58) and f8 (bit 61)rooksWhite = 0x0000000000000081 -- a1 (bit 0) and h1 (bit 7) , rooksBlack = 0x81000000000000 -- a8 (bit 56) and h8 (bit 63) queensWhite = 0x000000000000000 -- d1 (bit 3) queensBlack = 0x08000000000000 -- d8 (bit 59) , kingsWhite = 0x000000000000000 -- e1 (bit 4) , kingsBlack = 0x10000000000000 -- e8 (bit 60) }

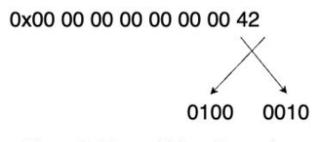


Figure 3: Hex to bitboard mapping

Bitboards

0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	
0	0	1	0	1	0	0	0	
0	0	0	0	0	0	0	0	
0	0	0	1	0	0	0	0	

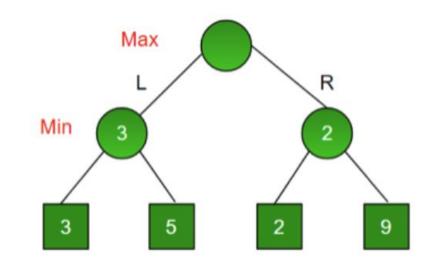
(a) Occupancy	(stop	before)	
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0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	
0	0	0	0	1	0	0	0	
1	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	1	
0	0	0	0	0	0	0	0	
0	0	0	0	1	0	0	0	
0	0	0	0	0	0	0	0	
(b) Captureable (stop on)								

1	0	0	0	0	0	0	0	
0	1	0	0	0	0	0	0	
0	0	1	0	0	0	0	0	
1	0	0	1	0	0	0	0	
1	0	0	0	1	0	0	1	
1	0	0	0	0	1	0	1	
1	1	0	0	0	0	1	1	
0	1	1	0	1	1	1	0	
(c) Queen moves from a1, h1								

Minimax Algorithm with optimizations

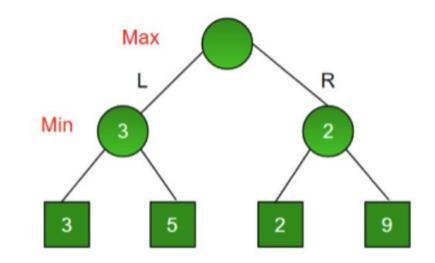
- Each turn the maximizing and min player switch roles and choose the most optimal branch
- Assumes each player plays optimally
- Space complexity 2ⁿ



Minimax Algorithm with optimizations

Caching

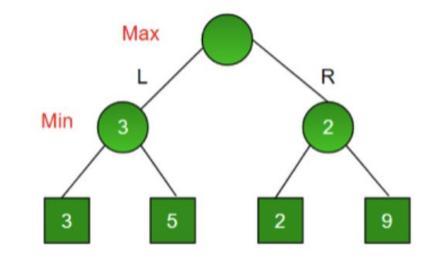
- Avoid recomputing expensive bitboard operations and comparisons
- Significant speedup at higher depths



Minimax Algorithm with optimizations

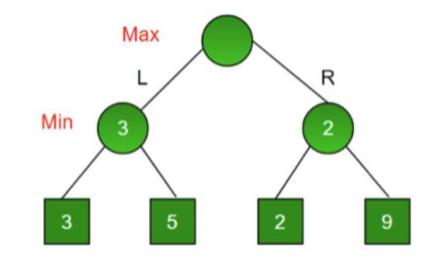
Pruning

- Avoid computing branches that we know the algorithm will never reach to save computational resources
- Has much more overhead than just caching but it takes runtime down even more aggressively

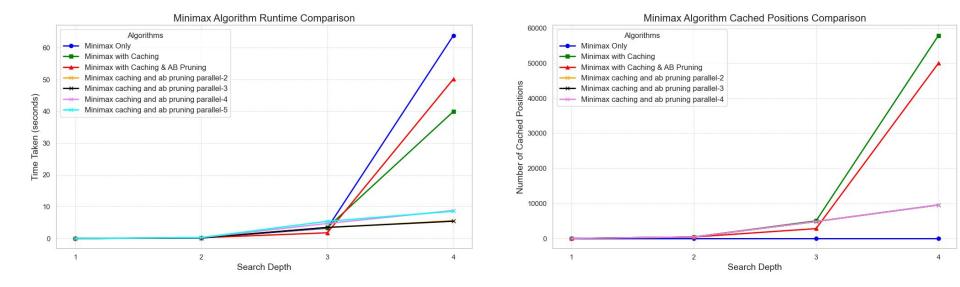


Parallelization

- Parallelize each top level minimax operation
- Leaves enough single threaded work, and breaks down the big work into sizeable chunks to take advantage of overhead
- Danger of exhausting system memory (24gb M3)



Preliminary Results



Preliminary Results - potential issues

- Exhausting resources
- Timing may not be fully accurate

